# Equivalences in PPL

October 15, 2013

Throughout, assume  $\alpha$  and  $\beta$  are (possibly complex) predicates. We'll use '= to denote logical equivalence.

### **Quantifier Exchange:**

$$\sim \exists \alpha \Rightarrow \models \forall \sim \alpha$$
$$\sim \forall \alpha \Rightarrow \models \exists \sim \alpha$$

### $\sim \exists \alpha \equiv \forall \sim \alpha$

Quantifier Replacement: 
$$\sim \exists \sim \alpha \implies \forall \alpha$$

$$\sim \exists \sim \alpha \Rightarrow \vdash \forall \alpha$$
$$\sim \forall \sim \alpha \Rightarrow \vdash \exists \alpha$$

### **Important Examples:**

### **Existence:**

$$\forall \alpha \vDash \exists \alpha, \mathbf{but}$$
:  $\forall (\alpha \supset \beta) \not\vDash \exists \alpha$ 

### **Conjunction:**

$$\exists (\alpha \& \beta) \not= \models \exists \alpha \& \exists \beta$$
$$\forall (\alpha \& \beta) = \models \forall \alpha \& \forall \beta$$

### Disjunction:

$$\exists (\alpha \lor \beta) \Rightarrow \vdash \exists \alpha \lor \exists \beta$$
$$\forall (\alpha \lor \beta) \Rightarrow \not \vdash \forall \alpha \lor \forall \beta$$

#### **Conditional:**

$$\forall (\alpha \supset \beta) \not= \models \forall \alpha \supset \forall \beta$$
$$\exists (\alpha \supset \beta) = \models \forall \alpha \supset \exists \beta$$

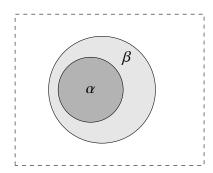
(You can work this out if you remember that  $\lceil \alpha \rceil \beta \rceil$  is equivalent to  $\lceil \sim \alpha \lor \beta \rceil$ , and use the rules given.)

### **Common Translation Schemes from English into PPL:**

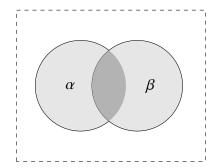
Some  $\alpha$ s are  $\beta$ s  $\Rightarrow$   $\exists (\alpha \& \beta)$ There are some  $\alpha$ s that are  $\beta$   $\Rightarrow$   $\exists (\alpha \& \beta)$ All  $\alpha$ s are  $\beta$ s  $\Rightarrow$   $\forall (\alpha \supset \beta)$ Every  $\alpha$  is a  $\beta$   $\Rightarrow$   $\forall (\alpha \supset \beta)$ No  $\alpha$ s are  $\beta$ s  $\Rightarrow$   $\neg \exists (\alpha \& \beta)$ No  $\alpha$ s are  $\beta$ s  $\Rightarrow$   $\neg \exists (\alpha \& \beta)$ No  $\alpha$ s are  $\beta$ s  $\Rightarrow$   $\forall (\alpha \supset \neg \beta)$ The  $\alpha$ s are exactly the  $\beta$ s  $\Rightarrow$   $\forall (\alpha \leftrightarrow \beta)$ 

# Venn Diagram Examples:

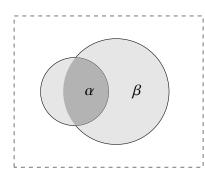




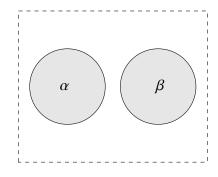
## $\exists (\alpha \& \beta)$ :



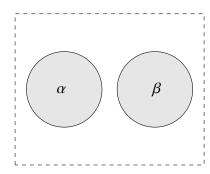
 $\sim \forall (\alpha \supset \beta)$ :



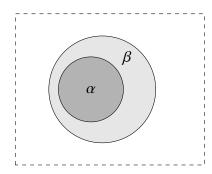
 $\sim \exists (\alpha \& \beta)$ :



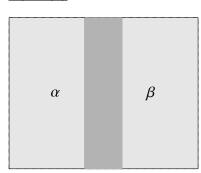
 $\forall (\alpha \supset \sim \beta)$ :



 $\underline{\sim}\exists(\alpha \& \sim\beta)$ :



 $\underline{\forall (\alpha \vee \beta)}$ :



 $\exists (\alpha \lor \beta)$ :

